

# Eigenvalue Methods for Sparse Tropical Polynomial Systems

Marianne Akian, Antoine Béréau, Stéphane Gaubert

CMAP, CNRS, École polytechnique, Institut Polytechnique de Paris, Inria

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- Given system of tropical polynomial equations or inequations, how to check the existence of, and then compute a solution in  $\mathbb{R}^n$ .
- Main tools in the classical setting include the theory of resultants, Macaulay matrices and effective Null- and Positivstellensatz.
- In this talk, we develop the tropical analog of the sparse Null- and Positivstellensatz, and explore the solvability of tropical polynomial systems by means of mean payoff games.

- 1 Tropical algebra and tropical polynomials**
- 2 The tropical Null- and Positivstellensatz**
- 3 Mean payoff games and tropical linear systems**

# I - Tropical algebra and tropical polynomials

- **Tropical semiring**  $\mathbb{T} := (\mathbb{R} \cup \{-\infty\}, \oplus, \odot, \mathbb{0}, \mathbb{1})$  with
  - ◇ addition  $\oplus := \max$ ;
  - ◇ multiplication  $\odot := +$ ;
  - ◇ zero element  $\mathbb{0} := -\infty$ ;
  - ◇ unit element  $\mathbb{1} := 0$ .
- Satisfies the usual properties of a field except **no additive inverse**.
- Tropical operations can be extended to vectors and matrices with coefficients in  $\mathbb{R} \cup \{-\infty\}$  to perform **tropical linear algebra**.

- A **formal tropical polynomial**  $p$  in  $n$  variables is a map

$$\begin{aligned} \mathbb{Z}^n &\longrightarrow \mathbb{R} \cup \{-\infty\} \\ \alpha &\longmapsto p_\alpha \end{aligned}$$

such that  $p_\alpha \neq \mathbb{0}$  for finitely many  $\alpha \in \mathbb{Z}^n$ . We denote  $p = \bigoplus_{\alpha \in \mathbb{Z}^n} p_\alpha X^\alpha$ .

- **Support** of  $p$ :  $\text{supp}(p) := \{\alpha \in \mathbb{Z}^n : p_\alpha \neq \mathbb{0}\}$ .
- **Polynomial function** associated to  $p$ :

$$\hat{p} : \begin{cases} \mathbb{R}^n &\longrightarrow \mathbb{R} \cup \{-\infty\} \\ x &\longmapsto \hat{p}(x) := \max_{\alpha \in \mathcal{A}} (p_\alpha + \langle x, \alpha \rangle) \end{cases}$$

with  $\mathcal{A} = \text{supp}(p)$ .

A point  $x \in (\mathbb{R} \cup \{-\infty\})^n$  is a **root** of a polynomial  $p$  whenever the maximum in the expression

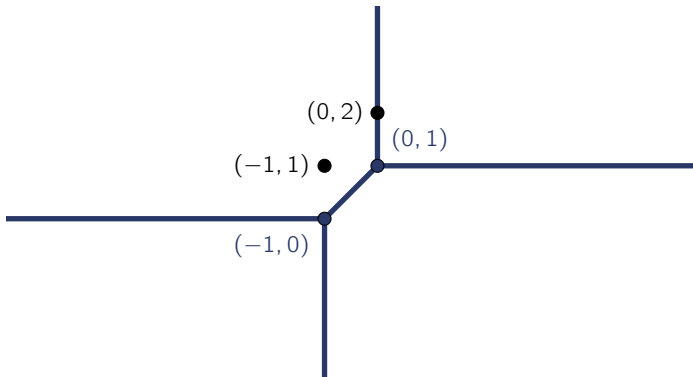
$$\hat{p}(x) = \bigoplus_{\alpha \in \mathcal{A}} p_{\alpha} \odot x^{\odot \alpha} = \max_{\alpha \in \mathcal{A}} (p_{\alpha} + \langle x, \alpha \rangle)$$

is attained for **at least two distinct values** of  $\alpha$ . This is denoted as  $p(x) \nabla \mathbb{0}$ .

**Example :** Let  $f_1 = 1 \oplus 2x_1 \oplus 1x_2 \oplus 1x_1x_2$ , then:

- $(0, 2)$  is a root of  $f_1$  since the maximum of  $\hat{f}_1(0, 2) = 3$  is attained simultaneously by the monomials  $1x_2$  and  $1x_1x_2$ ;
- $(-1, 1)$  is not a root of  $f_1$  since the maximum  $\hat{f}_1(-1, 1) = 2$  is attained *only* by the monomial  $1x_2$ .

The tropical hypersurface associated to the polynomial  $f_1$  is:





Likewise,  $y \in (\mathbb{R} \cup \{-\infty\})^m$  is said to be in the **tropical right null space** or **kernel** of a  $\ell \times m$  matrix  $A = (a_{ij})$  whenever for all  $1 \leq i \leq \ell$ , the maximum in the expression

$$\bigoplus_{j=1}^m a_{ij} \odot y_j = \max_{1 \leq j \leq m} (a_{ij} + y_j)$$

is achieved at least twice. This is also denoted as  $A \odot y \nabla \mathbb{0}$ .

*More on tropical geometry:* D. Maclagan and B. Sturmfels. *Introduction to Tropical Geometry*. Graduate Studies in Mathematics. American Mathematical Society, 2015.

## **II - The tropical Null- and Positivstellensatz**

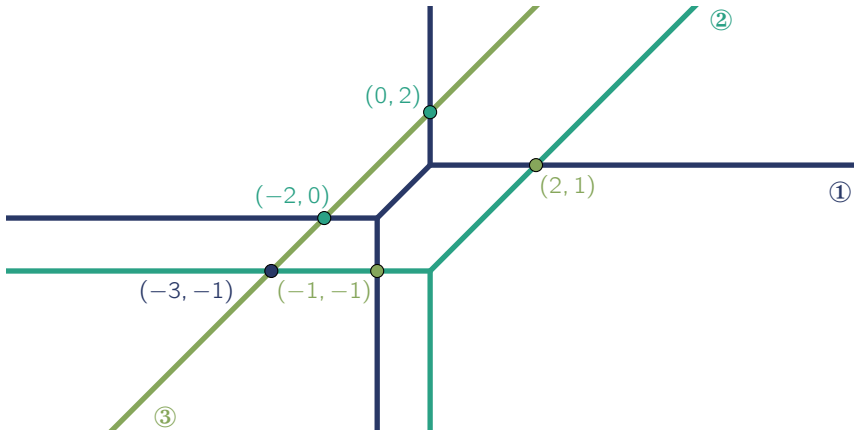
In the following, we fix a collection  $f = (f_1, \dots, f_k)$  of  $k$  formal tropical polynomials in  $n$  variables, with respective supports  $\mathcal{A} = (\mathcal{A}_1, \dots, \mathcal{A}_k)$  and degrees  $(d_1, \dots, d_k)$ .

**Problem:** Decide whether there is a common tropical zero  $x \in \mathbb{R}^n$ , that is such that  $f_i(x) \nabla 0$  for all  $1 \leq i \leq k$ .

**Remark:** The same question for a solution in  $(\mathbb{R} \cup \{-\infty\})^n$  reduces to the  $\mathbb{R}^n$  case by looking at all possible supports.

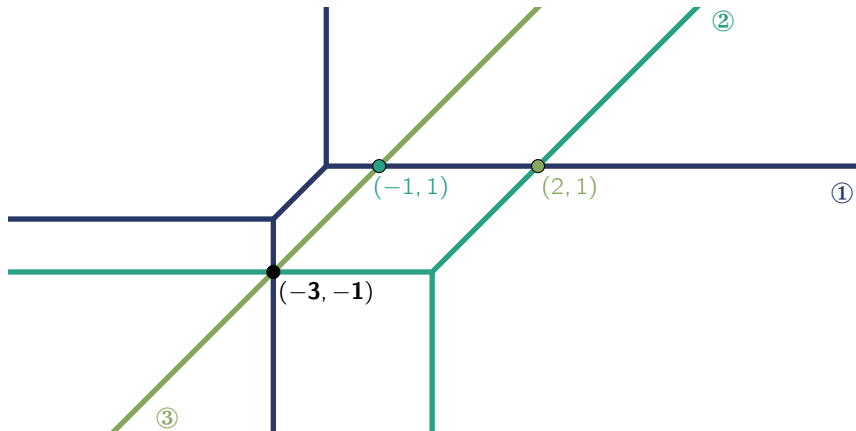
**Figure:** The arrangement of tropical varieties of the polynomials from the system

$$(E_1) : \begin{cases} f_1 = 1 \oplus 2x_1 \oplus 1x_2 \oplus 1x_1x_2 \\ f_2 = 0 \oplus 0x_1 \oplus 1x_2 \\ f_3 = 2x_1 \oplus 0x_2 \end{cases} .$$



**Figure:** The arrangement of tropical varieties of the polynomials from the system

$$(E_2) : \begin{cases} f_1 = 1 \oplus 4x_1 \oplus 1x_2 \oplus 3x_1x_2 \\ f_2 = 0 \oplus 0x_1 \oplus 1x_2 \\ f_3 = 2x_1 \oplus 0x_2 \end{cases} .$$



Link with classical varieties:

- Kapranov's theorem
- The Fundamental Theorem of Tropical Algebraic Geometry

Varied applications:

- celestial mechanics (Hampton, Moeckel)
- max-out networks (Montúfar, Ren, Zhang)
- chemical reaction networks (Dickenstein, Feliu, Radulescu, Shiu)
- emergency call center (Akian, Boyer, Gaubert)

The **Macaulay matrix** associated to  $f$  is the infinite matrix  $\mathcal{M} = (m_{(i,\alpha),\beta})$  indexed by  $([n] \times \mathbb{Z}^n) \times \mathbb{Z}^n$ , where  $m_{(i,\alpha),\beta}$  corresponds to the coefficient of  $X^\beta$  in the polynomial  $X^\alpha f_i$ .

$$\mathcal{M} = \begin{matrix} & 1 & x_1 & \cdots & x^\beta & \cdots \\ f_1 & * & * & \cdots & * & \cdots \\ x_1 f_1 & * & * & \cdots & * & \cdots \\ \vdots & \vdots & \vdots & \ddots & \vdots & \\ x^\alpha f_i & * & * & \cdots & f_{i,\beta-\alpha} & \cdots \\ \vdots & \vdots & \vdots & & \vdots & \ddots \end{matrix}$$

One approach to the **Nullstellensatz** is a linearization method reducing the search of a solution of a polynomial system to the search of nonzero elements in the kernel of the Macaulay matrix.

- The Macaulay matrix associated to  $f$  is the **infinite** matrix  $\mathcal{M} = (m_{(i,\alpha),\beta})$  indexed by  $([n] \times \mathbb{Z}^n) \times \mathbb{Z}^n$ , where  $m_{(i,\alpha),\beta}$  corresponds to the coefficient of  $X^\beta$  in the polynomial  $X^\alpha f_i$ .
- A finite subset  $\mathcal{E}$  of  $\mathbb{Z}^n$  yields a **finite** submatrix  $\mathcal{M}_{\mathcal{E}}$  of  $\mathcal{M}$  obtained by taking only the rows whose support is included in  $\mathcal{E}$  and the columns indexed by  $\mathcal{E}$ .
- Set  $\mathcal{M}_N := \mathcal{M}_{\mathcal{E}}$  for  $\mathcal{E} = \{\alpha \in \mathbb{N}^n : \alpha_1 + \cdots + \alpha_n \leq N\}$ .



**Conjecture [Grigoriev (2012)]:** There exists an integer  $N$  such that

$$\exists x \in \mathbb{R}^n \text{ such that } f_i(x) \nabla 0 \text{ for } i = 1, \dots, k$$

$$\iff$$

$$\exists y \in \mathbb{R}^m \text{ such that } \mathcal{M}_N \odot y \nabla 0 \text{ with } m = \binom{N+n}{n} .$$

Answer:

- **Grigoriev, Podolskii (2018):** true for

$$N = (n + 2)(d_1 + \dots + d_k) .$$

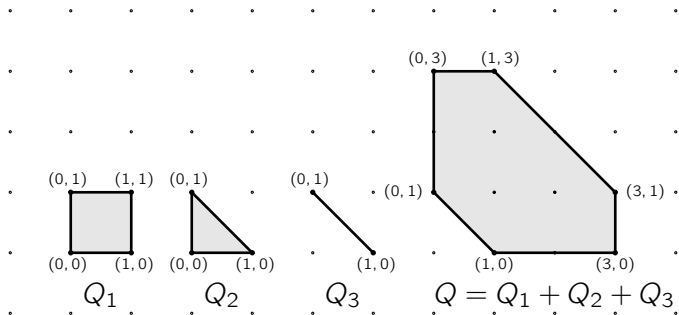
- **Akian, B., Gaubert (2023):** true for

$$N = d_1 + \dots + d_k - 1$$

(and even  $N = d_1 + \dots + d_k - n$  in most cases) + adapted approach for the case of sparse polynomials.

- For  $1 \leq i \leq k$ ,  $Q_i := \text{conv}(\mathcal{A}_i)$  is the **Newton polytope** of  $f_i$ .

**Example:** The Newton polytopes associated to both system  $(E_1)$  and system  $(E_2)$  and their Minkowski sum are as follow.



- **Canny-Emiris set** associated to  $f$ :  $\mathcal{E} = (Q + \delta) \cap \mathbb{Z}^n$  with  $\delta$  a generic vector in the linear space directing the affine hull of  $Q$ .

**Example:** Considering again the systems  $(E_1)$  and  $(E_2)$ , for

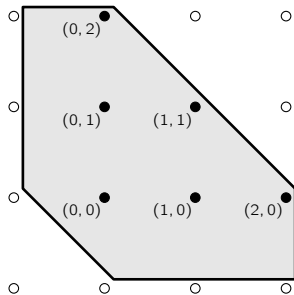
$$\delta = (-1 + \varepsilon, -1 + \varepsilon)$$

with  $\varepsilon > 0$  sufficiently small, we obtain the Canny-Emiris set

$$\mathcal{E} := (Q + \delta) \cap \mathbb{Z}^n = \{(0, 0), (1, 0), (0, 1), (2, 0), (1, 1), (0, 2)\}$$

corresponding to the set of monomials  $\{1, x_1, x_2, x_1^2, x_1x_2, x_2^2\}$ .

**Figure:** The polytope  $Q + \delta$  with  $\delta = (-0.9, -0.9)$ .



## Nullstellensatz for Sparse Tropical Polynomial Systems

The system  $f \nabla \mathbb{0}$  has a solution  $x \in \mathbb{R}^n$  iff there exists a vector  $y \in \mathbb{R}^{\mathcal{E}'}$  in the tropical right null space of the submatrix  $\mathcal{M}_{\mathcal{E}'}$  of  $\mathcal{M}$ , where  $\mathcal{E}'$  is any subset of  $\mathbb{Z}^n$  containing a nonempty Canny-Emiris set  $\mathcal{E}$ .

**Corollary:** The system  $f \nabla \mathbb{0}$  has a solution  $x \in \mathbb{R}^n$  if and only if the truncated Macaulay tropical linear system  $\mathcal{M}_N \odot y \nabla \mathbb{0}$  has a solution  $y \in \mathbb{R}^m$  for

$$N = d_1 + \cdots + d_k - 1 ,$$

where  $d_i = \deg(f_i)$  for all  $1 \leq i \leq k$ . Moreover, if  $Q$  has full dimension, then one can take  $N = d_1 + \cdots + d_k - n$  in the previous statement.

**Example:** The matrix associated with system  $(E_1)$  is

$$\mathcal{M}_{\mathcal{E}}^{(1)} = \begin{matrix} & & & 1 & x_1 & x_2 & x_1^2 & x_1x_2 & x_2^2 \\ f_1 & & & 1 & 2 & 1 & & 1 & \\ f_2 & & & 0 & 0 & 1 & & & \\ x_1f_2 & & & & 0 & & 0 & 1 & \\ x_2f_2 & & & & & 0 & & 0 & 1 \\ f_3 & & & & 2 & 0 & & & \\ x_1f_3 & & & & & & 2 & 0 & \\ x_2f_3 & & & & & & & 2 & 0 \end{matrix} .$$

There is no finite vector in its tropical right null space and thus there is no finite solution to the equation  $f \nabla 0$ .

**Example:** The matrix associated with system  $(E_2)$  is

$$\mathcal{M}_{\mathcal{E}}^{(2)} = \begin{matrix} & & & 1 & x_1 & x_2 & x_1^2 & x_1x_2 & x_2^2 \\ f_1 & & & 1 & 4 & 1 & & 3 & \\ f_2 & & & 0 & 0 & 1 & & & \\ x_1f_2 & & & & 0 & & 0 & 1 & \\ x_2f_2 & & & & & 0 & & 0 & 1 \\ f_3 & & & & 2 & 0 & & & \\ x_1f_3 & & & & & & 2 & 0 & \\ x_2f_3 & & & & & & & 2 & 0 \end{matrix}.$$

The vector  $y = \text{ver}(-3, -1) = (0, -3, -1, -6, -4, -2)$  is finite and is in the tropical right null space of the previous matrix, hence there is a finite solution to the equation  $f \nabla 0$ , which is indeed given by  $(-3, -1)$ .

- Let  $f^\pm = (f_1^\pm, \dots, f_k^\pm)$  be two collections of tropical polynomials. For  $1 \leq i \leq k$ , denote by  $\mathcal{A}_i^\pm$  the support of  $f_i^\pm$ .
- Set  $\triangleright = (\triangleright_1, \dots, \triangleright_k)$  a collection of relations, with  $\triangleright_i \in \{\geq, =, >\}$  for  $1 \leq i \leq k$ .

We denote by  $f^+(x) \triangleright f^-(x)$  the system

$$\max_{\alpha \in \mathcal{A}_i^+} (f_{i,\alpha}^+ + \langle \alpha, x \rangle) \triangleright_i \max_{\alpha \in \mathcal{A}_i^-} (f_{i,\alpha}^- + \langle \alpha, x \rangle) \text{ for all } 1 \leq i \leq k$$

of unknown  $x \in (\mathbb{R} \cup \{-\infty\})^n$ .



- Let  $\mathcal{M}^\pm$  be the Macaulay matrices associated to  $f^\pm$  – i.e. with entries  $f_{i,\beta}^\pm$ . For any subset  $\mathcal{E}$  of  $\mathbb{Z}^n$ , denote by  $\mathcal{M}_\mathcal{E}^\pm$  the submatrices of  $\mathcal{M}^\pm$  by taking only the row indices  $(i, \alpha) \in [k] \times \mathbb{Z}^n$  such that the supports of the rows  $(i, \alpha)$  of both  $\mathcal{M}^+$  and  $\mathcal{M}^-$  is included in  $\mathcal{E}$  and the column indices given by  $\mathcal{E}$ .
- Finally, denote by  $\mathcal{M}_\mathcal{E}^+ \odot y \triangleright \mathcal{M}_\mathcal{E}^- \odot y$  the following system of tropical linear inequalities:

$$\max_{\beta \in \mathcal{E}} \left( \mathcal{M}_{(i,\alpha),\beta}^+ + y_\beta \right) \triangleright_i \max_{\beta \in \mathcal{E}} \left( \mathcal{M}_{(i,\alpha),\beta}^- + y_\beta \right) \text{ for all } 1 \leq i \leq k.$$

Let  $\tilde{Q} = r_1 Q_1 + \dots + r_k Q_k$ , where  $Q_i = \text{conv}(\mathcal{A}_i^+ \cup \mathcal{A}_i^-)$  for  $i = 1, \dots, k$ , and

$$r_i = \begin{cases} \min(|\mathcal{A}_i^-|, n + 1) & \text{if } \triangleright_i \in \{\geq, >\} \\ \min(\max(|\mathcal{A}_i^-|, |\mathcal{A}_i^+|), n + 1) & \text{if } \triangleright_i \in \{=\} \end{cases} .$$

We now call **Canny-Emiris subsets** of  $\mathbb{Z}^n$  associated to the pair of collections  $(f^+, f^-)$  any set  $\mathcal{E}$  of the form

$$\mathcal{E} := (\tilde{Q} + \delta) \cap \mathbb{Z}^n ,$$

where  $\delta$  is a generic vector in  $V + \mathbb{Z}^n$ , with  $V$  the direction of the affine hull of  $\tilde{Q}$ .

## Tropical Positivstellensatz

There exists a solution  $x \in \mathbb{R}^n$  to the system  $f^+(x) \triangleright f^-(x)$  if and only if there exists a vector  $y \in \mathbb{R}^{\mathcal{E}'}$  satisfying  $\mathcal{M}_{\mathcal{E}'}^+ \odot y \triangleright \mathcal{M}_{\mathcal{E}'}^- \odot y$ , where  $\mathcal{E}'$  is any subset of  $\mathbb{Z}^n$  containing a nonempty Canny-Emiris subset  $\mathcal{E}$  of  $\mathbb{Z}^n$  associated to the pair  $(f^+, f^-)$ .

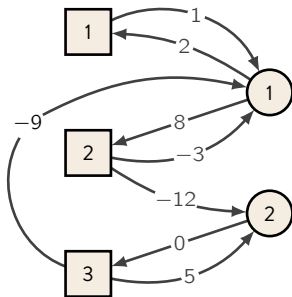
**Corollary:** The inclusion of basic tropical semialgebraic sets can be reduced to solving a set of tropical linear (in)equalities.

# III - Mean payoff games and tropical linear systems

Mean payoff games (See **Gillette (1957)**,  
**Gurvich, Karzanov, Khachiyan (1988)**,  
**Zwick, Patterson (1996)**):

- $G = (I \sqcup J, E)$  a (finite) oriented weighted bipartite graph;
- game with two players Min and Max: each turn, from the current state  $i \in I$ , player Max chooses a state  $j \in J$  such that  $(i, j)$  is an arc of  $G$  with weight  $b_{ij}$  and obtains a payment of  $b_{ij}$  from player Min, then player Min from state  $j \in J$ , chooses the next state  $k \in I$  along an arc  $(k, j)$  with weight  $a_{kj}$ , and receives in turn a payment of  $a_{kj}$  from player Max;
- the winner is the player who gets the highest average payment per turn;
- set  $A = (a_{ij})_{(i,j) \in I \times J}$  et  $B = (b_{ij})_{(i,j) \in I \times J}$ .

**Example :** Let  $G$  be the following graph:

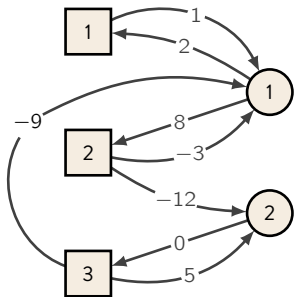


One has  $A = \begin{pmatrix} 2 & -\infty \\ 8 & -\infty \\ -\infty & 0 \end{pmatrix}$  and  $B = \begin{pmatrix} 1 & -\infty \\ -3 & -12 \\ -9 & 5 \end{pmatrix}$ .

**Theorem [Akian, Gaubert, Guterman (2012)] :** For all  $j \in J$ , player Max has a winning positional strategy for the *mean pay-off game* given by the payment matrices  $A$  and  $B$  by playing the **initial move  $j$**  iff there exists a solution  $y \in (\mathbb{R} \cup \{-\infty\})^J$  of the tropical matrix inequality  $A \odot y \leq B \odot y$  such that  $y_j \neq \mathbb{0}$ .

The winning initial moves correspond to the **support** of the solutions of the inequality  $A \odot y \leq B \odot y$ .

In the previous example,



$$\text{one has } A \odot y \leq B \odot y \iff \begin{cases} 2 + y_1 \leq 1 + y_1 \\ 8 + y_1 \leq \max(-3 + y_1, -12 + y_2) \\ y_2 \leq \max(-9 + y_1, 5 + y_2). \end{cases}$$

The first inequality shows that every solution  $y \in (\mathbb{R} \cup \{-\infty\})^2$  must satisfy  $y_1 = 0$ , which implies that the two other inequalities are satisfied for all values of  $y_2 \in \mathbb{R} \cup \{-\infty\}$ .

This translates into the fact that the **move 1** is a **losing move** for player Max, while the **move 2** is a **winning move**.



- **Shapley operator** associated to a mean payoff game

$$T : (\mathbb{R} \cup \{\pm\infty\})^J \longrightarrow (\mathbb{R} \cup \{\pm\infty\})^J$$

$$y = (y_j)_{j \in J} \longmapsto \left( \min_{i \in I} -a_{ij} + \left( \max_{k \in J} b_{ik} + y_k \right) \right)_{j \in J}$$

- **value** of the game:  $\chi(T) = \lim_{n \rightarrow +\infty} \frac{T^n(0)}{n}$

**Corollary:**  $\exists y \in \mathbb{R}^J$  such that  $A \odot y \leq B \odot y$  iff  $\min_{j \in J} \chi_j(T) \geq 0$ .

Link with nonlinear eigenvalue theory:

$$\begin{aligned} & \min\{\chi_j(T) : j \in J\} \\ &= \sup\{\lambda \in \mathbb{R} : \exists u \in \mathbb{R}^J, T(u) \geq \lambda + u\} \\ &= \inf\{\lambda \in \mathbb{R} \cup \{+\infty\} : \exists u \in (\mathbb{R} \cup \{+\infty\})^J, u \neq +\infty, T(u) \leq \lambda + u\} \\ &= \inf\{\lambda \in \mathbb{R} \cup \{+\infty\} : \exists u \in (\mathbb{R} \cup \{+\infty\})^J, u \neq +\infty, T(u) = \lambda + u\} . \end{aligned}$$

In particular  $\chi(T) \equiv \lambda \in \mathbb{R}$  iff the nonlinear eigenproblem

$$T(u) = \lambda + u$$

has a solution  $(\lambda, u) \in \mathbb{R} \times \mathbb{R}^J$ .

The existence of a polynomial time algorithm to solve mean payoff games is an open problem since 1988, but there exist practically fast methods (value/policy iteration algorithms).

For a Shapley operator  $T : (\mathbb{R} \cup \{+\infty\})^J \rightarrow (\mathbb{R} \cup \{+\infty\})^J$ , define the Krasnoselskii-Mann damped Shapley operator  $T_{KM}$  by  $T_{KM}(u) = \frac{u+T(u)}{2}$  for all  $u \in (\mathbb{R} \cup \{+\infty\})^J$ . Then  $\chi(T_{KM}) = \frac{\chi(T)}{2}$

We propose the following value iteration algorithm.

# Value iteration algorithm

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## Algorithm 1: Value iteration algorithm with widening.

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**input:**  $T$  a Shapley operator from  $(\mathbb{R} \cup \{+\infty\})^J$  to  $(\mathbb{R} \cup \{+\infty\})^J$   
 $\varepsilon > 0$  the approximation error for comparisons  
 $N^*$  a timeout on the number of iterations which guarantees the existence of a solution whenever reached

**output:** Decides the feasibility of the system  $y \leq T(y)$  in  $\mathbb{R}^J$

```
1 initialization
2  $u := 0 \in \mathbb{R}^J$ 
3  $v := 0 \in \mathbb{R}^J$ 
4  $N := 0$ 
5 repeat
6     /* Value iteration step */
7      $u := v$ 
8      $v := u \wedge T(u)$ 
9      $N := N + 1$ 
10    /* Widening step */
11     $I := \{i : v_i \geq -\varepsilon + u_i\}$ 
12     $\hat{u} := (\hat{u}_i) \in (\mathbb{R} \cup \{+\infty\})^m$  with  $\begin{cases} \hat{u}_i = +\infty & \text{if } i \in I \\ \hat{u}_i = u_i & \text{otherwise} \end{cases}$ 
13     $\hat{v} := T(\hat{u})$ 
14 until  $v \geq -\varepsilon + u$  or  $v \ll -\varepsilon + u$  or  $\hat{v} \ll -\varepsilon + \hat{u}$  or  $\min_{i \in J} (u_i) < -(|J| - 1)W$  or  $N \geq N^*$ 
15 if  $v \ll -\varepsilon + u$  or  $\hat{v} \ll -\varepsilon + \hat{u}$  or  $\min_{i \in J} (u_i) < -(|J| - 1)W$  then
16     return "Unfeasible"
17 else
18     return "Feasible"
```

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### Correction and termination of the value iteration algorithm

Algorithm 1 applied to the Krasnoselskii-Mann damped Shapley operator  $T_{KM}$  correctly decides (in exact arithmetic) the feasibility of a tropical linear system with integer coefficients in  $N^* = \mathcal{O}(|J|^2W)$  iterations for  $\varepsilon < \frac{1}{\min(|I|, |J|)}$ , where  $W$  is an upper bound on the maximal non  $-\infty$  coefficients of  $A$  and  $B$ .

- Algorithm 1 is also correct and terminates in approximate arithmetics for sufficiently small approximation errors.
- Since the cost of each evaluation of the operator  $T$  is pseudo-polynomial, Algorithm 1 is in pseudo-polynomial complexity.
- To be compared with policy iteration algorithms.

Let  $f^\pm = (f_1^\pm, \dots, f_k^\pm)$  be two collections of tropical polynomials and let  $d = \max_{1 \leq i \leq k} \deg(f_i^\pm)$  and  $W = \max_{1 \leq i \leq k} \|f_i^\pm\|_\infty$ , and for  $\epsilon \in \{\pm 1\}^n$ , denote by  $\epsilon \mathbb{R}_{\geq 0}^n$  the orthant  $\{x \in \mathbb{R}^n : \epsilon_j x_j \geq 0 \text{ for all } 1 \leq j \leq n\}$ . Then:

### Short model property

- The vertices of every polyhedral complex  $\{x \in \mathbb{R}^n : f_i^+(x) \geq f_i^-(x)\} \cap \epsilon \mathbb{R}_{\geq 0}^n$  are included in a  $\|\cdot\|_\infty$ -ball of radius  $2n(2d)^{n-1}W$  centered at point 0
- Moreover, if all the coefficients of the polynomials  $f_i^\pm$  are integer, these vertices have coordinates that are rational numbers with a denominator bounded above by  $(2d)^n$ .

## Dichotomic search method

- Solve the system

$$\begin{cases} f^+(x) \triangleright f^-(x) \\ a \leq x_1 \leq b \end{cases}$$

for varying values of  $a$  and  $b$ .

- If  $|b - a| < \frac{1}{(2d)^n}$  then one can deduce the first coordinate of a solution.
- Fix the value of  $x_1$  and repeat with  $x_2, \dots, x_n$ .

The dichotomic search method returns a rational solution of this system (or decides that there is none) in  $\mathcal{O}(\log(n(2d)^{2n-1}W))$  calls to a weak mean payoff oracle.



## Path-following method

- Solve the system

$$f^+(\zeta, x_2, \dots, x_n) \triangleright f^-(\zeta, x_2, \dots, x_n)$$

for varying values of  $\zeta$ .

- Linearize the above system and consider the associated mean payoff game with its shapley operator  $T_\zeta$ .
- $\phi : \zeta \mapsto \min_{j \in J} \chi_j(T_\zeta)$  is a continuous, Lipschitz, piecewise affine function.
- Computing  $\phi$  with a pivoting algorithm yields the projection of the solution set onto the first coordinate.

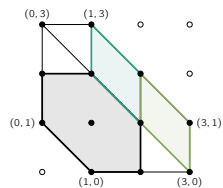
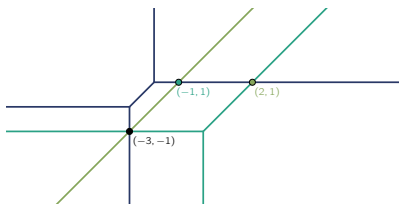
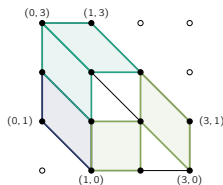
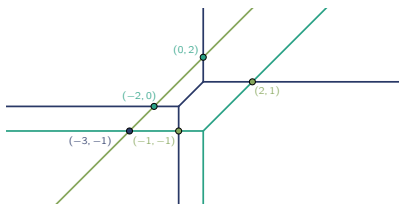
Python implementation of the algorithm available at:

<https://gitlab.inria.fr/abereau/tropical-polynomial-system-solving>

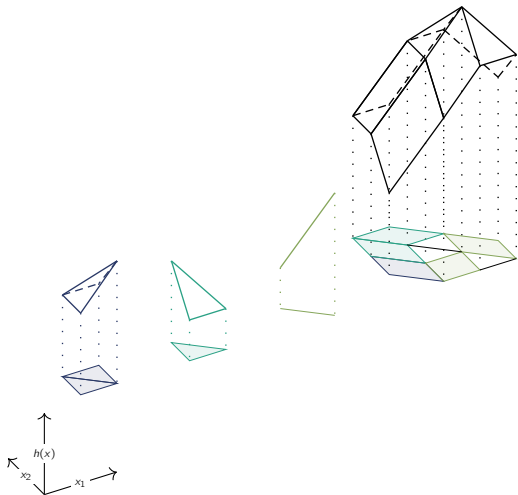
The bottleneck resides mainly in the computation of the Minkowski sum of the Newton polytopes of the polynomials of the system.

**Open problem:** Can the degree bound be improved in the Positivstellensatz (no tight example found yet)?

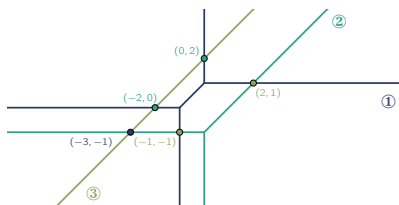
# Thank you for your attention!



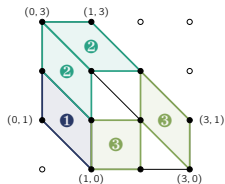
**Figure:** The subdivision of  $Q$  associated to  $(E_1)$  arises from the projection of the Minkowski sum of the hypographs of the lifted Newton polytopes.



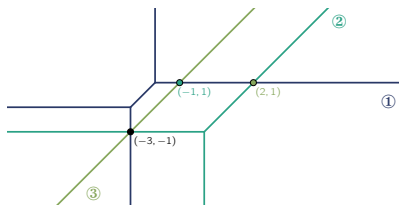
(a) The arrangement of tropical varieties of the polynomials from the system  $(E_1)$ .



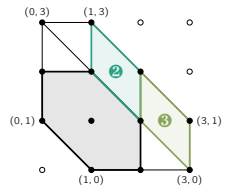
(b) The subdivision of  $Q$  associated to  $(E_1)$ .



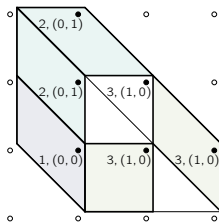
(c) The arrangement of tropical varieties of the polynomials from the system  $(E_2)$ .



(d) The subdivision of  $Q$  associated to  $(E_2)$ .



**Figure:** The polytope  $Q + \delta$ , with the integer points inside the maximal dimensional cells of the decomposition of  $Q + \delta$  labelled by the row content the cell they belong to.



This configuration yields the following nonsingular square submatrix of  $\mathcal{M}_{\mathcal{E}}^{(1)}$

$$\mathcal{M}_{\mathcal{E}\mathcal{E}}^{(1)} = \begin{matrix} & & & & 1 & x_1 & x_2 & x_1^2 & x_1 x_2 & x_2^2 \\ \begin{matrix} (0,0) \rightarrow f_1 \\ (1,0) \rightarrow f_3 \\ (0,1) \rightarrow f_2 \\ (2,0) \rightarrow x_1 f_3 \\ (1,1) \rightarrow x_2 f_3 \\ (0,2) \rightarrow x_2 f_2 \end{matrix} & \begin{pmatrix} 1 & 2 & 1 & & & & & & & \\ 0 & 0 & 1 & & & & & & & \\ & & & 2 & 0 & & & & & \\ & & & & 2 & 0 & & & & \\ & & & & 0 & 0 & & & & \\ & & & & & 0 & & & & 1 \end{pmatrix} & \end{matrix}.$$

## The Shapley-Folkman Lemma

Let  $A_1, \dots, A_k \subseteq \mathbb{R}^n$ , and let

$$x \in \sum_{i=1}^k \text{conv}(A_i) .$$

Then there is an index set  $I \subseteq \{1, \dots, k\}$  with  $|I| \leq n$  such that

$$x \in \sum_{i \in I} \text{conv}(A_i) + \sum_{i \in \{1, \dots, k\} \setminus I} A_i .$$

**Corollary:** If  $\sum_{i=1}^k \text{conv}(A_i)$  has (affine) dimension  $d < n$ , then the index set  $I$  can be chosen such that  $|I| \leq d$ .